# Introduction to OpenACC

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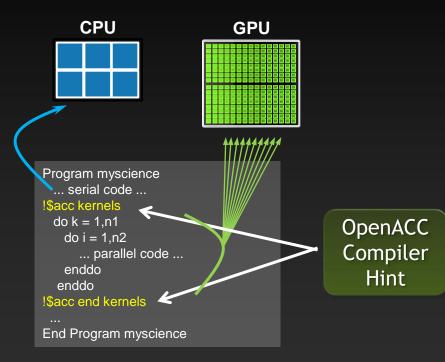
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# What is OpenACC?

It is a directive based standard to allow developers to take advantage of accelerators such as GPUs from NVIDIA and AMD, Intel's Xeon Phi, FPGAs, and even DSP chips.



# **Directives**



Simple compiler hints from coder.

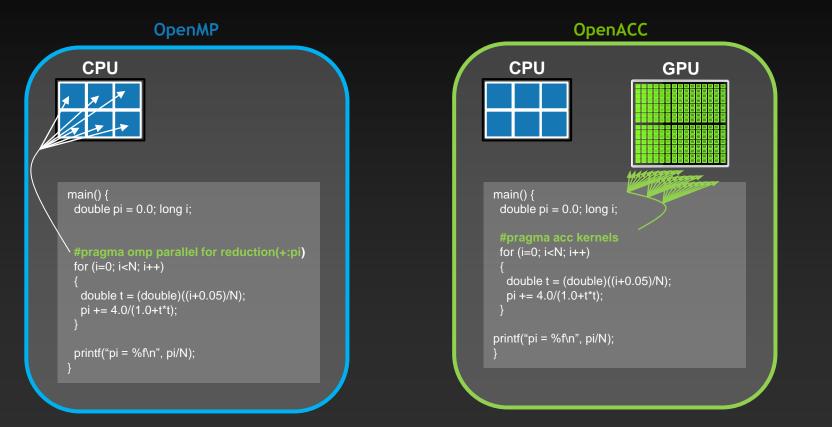
Compiler generates parallel threaded code.

Ignorant compiler just sees some comments.

Your original Fortran or C code



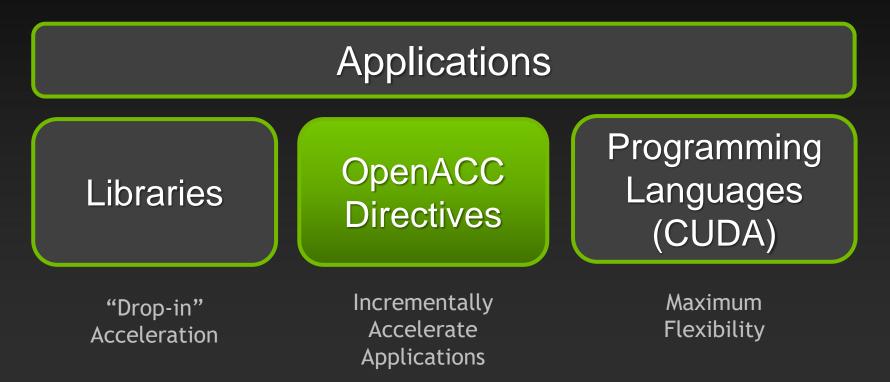
# Familiar to OpenMP Programmers



#### More on this later!



# How Else Would We Accelerate Applications?





# Key Advantages Of This Approach

- High-level. No involvement of OpenCL, CUDA, etc.
- Single source. No forking off a separate GPU code. Compile the same program for accelerators or serial; non-GPU programmers can play along.
- Efficient. Experience shows very favorable comparison to low-level implementations of same algorithms.
- Performance portable. Supports GPU accelerators and co-processors from multiple vendors, current and future versions.
- Incremental. Developers can port and tune parts of their application as resources and profiling dictates. No wholesale rewrite required. Which can be <u>quick.</u>



## **True Standard**

Full OpenACC specifications (now on 3.2) available online

http://www.openacc-standard.org

- Quick reference card also available and useful
- Implementations available now from PGI, Cray, CAPS and GCC.
- GCC version of OpenACC started in 5.x, but use at least 10.x
- Best free option is very probably PGI Community version: http://www.pgroup.com/products/community.htm

#### The OpenACC<sup>™</sup> API QUICK REFERENCE GUIDE

The OpenACC Application Program Interface describes a collection of compiler directives to specify loops and regions of code in standard C, C++ and Fortran to be offloaded from a host CPU to an attached accelerator, providing portability across operating systems, host CPUs and accelerators.

Most OpenACC directives apply to the immediately following structured block or loop; a structured block is a single statement or a compound statement (C or C++) or a sequence of statements (Fortan) with a single entry point at the top and a single exit at the bottom.





#### **OPENACC** Resources

Guides • Talks • Tutorials • Videos • Books • Spec • Code Samples • Teaching Materials • Events • Success Stories • Courses • Slack • Stack Overflow

## FREE Compilers





#### Resources https://www.openacc.org/resources



#### Compilers and Tools

https://www.openacc.org/tools



#### **Success Stories**

https://www.openacc.org/success-stories



# Events

The OpenACC Community organizes a variety of events throughout the year. Events vary from talks at conferences to workshops, haskathons, solline courses and User Group meetings. Join our events around the world to learn OpenACC programming and to participate in activities with the OpenaCC Lister Frown.



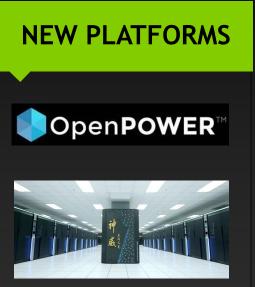
#### Hackathons

Hackabons are five day intensive hands on mentoring sessions. They are designed to help computational scientific part thair applications to 60% using literative, OpenRCC, CUDA and other tools. They are currently added by the OAR Helps Laedenhip Compositive filterative (COCR) and the Mellegia National Multicortory (ORA). The the Multi schedule and registration details pleases visit <u>https://www.add.oral.gov/laming.oven/2011.gov</u>.





# **Serious Adoption**



Sunway TaihuLight Built around OpenACC

#### GROWING COMMUNITY



- 6,000+ enabled developers
- Hackathons constantly
- Diverse online community

#### PORTING SUCCESS

- Five of 13 CAAR codes using OpenACC
- Gaussian ported to Tesla with OpenACC
- FLUENT using OpenACC in R18 production release



# **A Few Cases**

Designing circuits for quantum

computing

UIST, Macedonia

Reading DNA nucleotide sequences Shanghai JiaoTong University



#### 4 directives

16x faster

HydroC- Galaxy Formation
<u>PRACE Benchmark Code</u>, CAPS



1 week

3x faster



1 week

40x faster

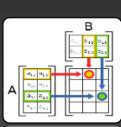
Real-time Derivative Valuation

Opel Blue, Ltd



Few hours

70x faster





6.4x faster



Extracting image features in realtime

Aselsan



3 directives

4.1x faster

Matrix Matrix Multiply

Independent Research Scientist

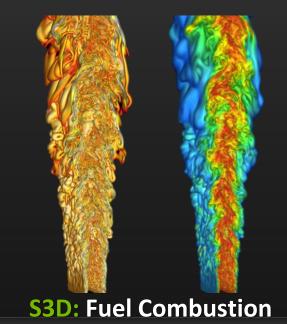
## **A Champion Case**

# **4x FasterJaguarTitan**42 days10 days

#### Modified <1% Lines of Code

15 PF! One of fastest simulations ever!

Design alternative fuels with up to 50% higher efficiency





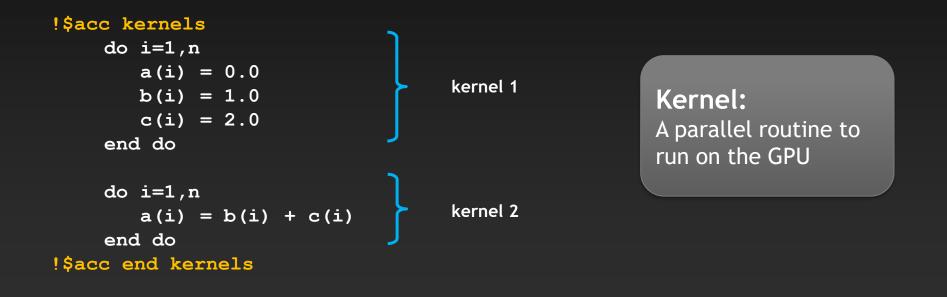
#### A Simple Example: SAXPY

#### SAXPY in C SAXPY in Fortran subroutine saxpy(n, a, x, y) void saxpy(int n, real :: x(:), y(:), a float a, integer :: n, i float \*x, !\$acc kernels float \*restrict y) do i=1.n y(i) = a\*x(i)+y(i)#pragma acc kernels enddo for (int i = 0; i < n; ++i) !\$acc end kernels y[i] = a\*x[i] + y[i];end subroutine saxpy } \$ From main program // Somewhere in main \$ call SAXPY on 1M elements // call SAXPY on 1M elements call saxpy(2\*\*20, 2.0, x\_d, y\_d) <u>saxpy(1<<20, 2.0, x, y);</u> . . .



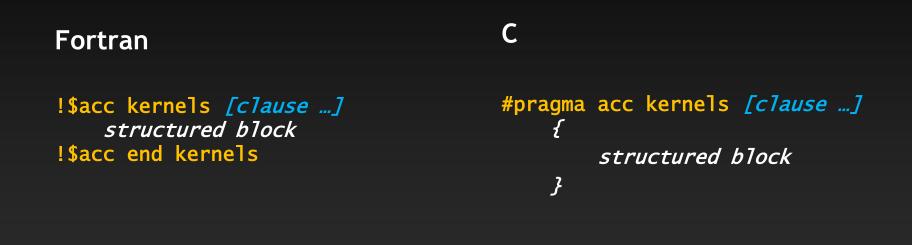
## kernels: Our first OpenACC Directive

We request that each loop execute as a separate *kernel* on the GPU. This is an incredibly powerful directive.





# **General Directive Syntax and Scope**

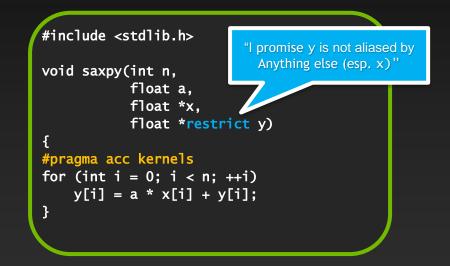


I may indent the directives at the natural code indentation level for readability. It is a common practice to always start them in the first column (ala #define/#ifdef). Either is fine with C or Fortran 90 compilers.



## **Complete SAXPY Example Code**

```
int main(int argc, char **argv)
 int N = 1 << 20; // 1 million floats
 if (argc > 1)
   N = atoi(argv[1]);
 float *x = (float*)malloc(N * sizeof(float));
 float *y = (float*)malloc(N * sizeof(float));
 for (int i = 0; i < N; ++i) {
   x[i] = 2.0f;
   y[i] = 1.0f;
  }
 saxpy(N, 3.0f, x, y);
  return 0;
```





## C Detail: the restrict keyword

- Standard C (as of C99).
- Important for optimization of serial as well as OpenACC and OpenMP code.
- Promise given by the programmer to the compiler for a pointer float \*restrict ptr

Meaning: "for the lifetime of ptr, only it or a value directly derived from it (such as ptr + 1) will be used to access the object to which it points"

- Limits the effects of pointer aliasing
- OpenACC compilers often require restrict to determine independence
  - Otherwise the compiler can't parallelize loops that access ptr
  - Note: if programmer violates the declaration, behavior is undefined

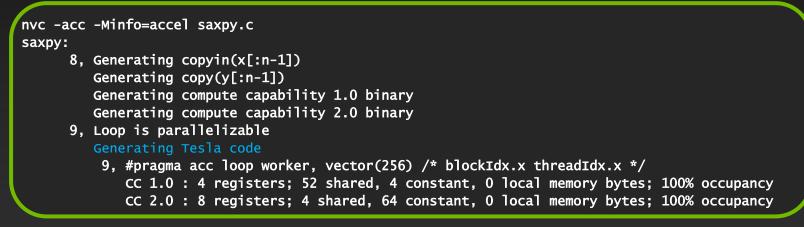


## **Compile and Run**

C: nvc -acc -Minfo=accel saxpy.c

Fortran: nvfortran -acc -Minfo=accel saxpy.f90

**Compiler Output** 



Run: a.out

#### Compare: Partial CUDA C SAXPY Code Just the subroutine

```
global void saxpy kernel( float a, float* x, float* y, int n ){
 int i;
 i = blockIdx.x*blockDim.x + threadIdx.x;
 if(i \le n) x[i] = a*x[i] + y[i];
void saxpy( float a, float* x, float* y, int n ){
  float *xd, *yd;
  cudaMalloc( (void**)&xd, n*sizeof(float) );
  cudaMalloc( (void**)&yd, n*sizeof(float) ); cudaMemcpy( xd, x, n*sizeof(float),
                     cudaMemcpyHostToDevice );
  cudaMemcpy( vd, v, n*sizeof(float),
                     cudaMemcpyHostToDevice );
  saxpy kernel<<< (n+31)/32, 32 >>>( a, xd, yd, n );
  cudaMemcpy( x, xd, n*sizeof(float),
                     cudaMemcpyDeviceToHost );
  cudaFree( xd ); cudaFree( yd );
```



#### Compare: Partial CUDA Fortran SAXPY Code Just the subroutine

```
module kmod
 use cudafor
contains
 attributes(global) subroutine saxpy kernel(A,X,Y,N)
  real(4), device :: A, X(N), Y(N)
  integer, value :: N
  integer :: i
  i = (blockidx%x-1)*blockdim%x + threadidx%x
  if(i \le N) X(i) = A*X(i) + Y(i)
 end subroutine
end module
 subroutine saxpy(A, X, Y, N)
  use kmod
  real(4) :: A, X(N), Y(N)
  integer :: N
  real(4), device, allocatable, dimension(:):: &
                 Xd, Yd
  allocate(Xd(N), Yd(N))
  Xd = X(1:N)
  Yd = Y(1:N)
  call saxpy kernel \langle \langle (N+31)/32, 32 \rangle \rangle \langle A, Xd, Yd, N \rangle
  X(1:N) = Xd
  deallocate(Xd, Yd)
 end subroutine
```



## Again: Complete SAXPY Example Code

#### Main Code

```
int main(int argc, char **argv)
 int N = 1 << 20; // 1 million floats
 if (argc > 1)
   N = atoi(argv[1]);
 float *x = (float*)malloc(N * sizeof(float));
  float *y = (float*)malloc(N * sizeof(float));
  for (int i = 0; i < N; ++i) {
   x[i] = 2.0f;
   y[i] = 1.0f;
  }
 saxpy(N, 3.0f, x, y);
  return 0;
```

#### **Entire Subroutine**

#include <stdlib.h>

```
void saxpy(int n,
float a,
float *x,
float *restrict y)
```

```
#pragma acc kernels
for (int i = 0; i < n; ++i)
    y[i] = a * x[i] + y[i];</pre>
```



# **Big Difference!**

- With CUDA, we changed the structure of the old code. Non-CUDA programmers can't understand new code. It is not even ANSI standard code.
- We have separate sections for the host code and the GPU code. Different flow of code. Serial path now gone forever.
- Where did these "32"s and other mystery numbers come from? This is a clue that we have some hardware details to deal with here.
- Exact same situation as assembly used to be. How much hand-assembled code is still being written in HPC now that compilers have gotten so efficient?



#### This looks easy! Too easy...

- If it is this simple, why don't we just throw kernel in front of every loop?
- Better yet, why doesn't the compiler do this for me?

The answer is that there are two general issues that prevent the compiler from being able to just automatically parallelize every loop.

- Data Dependencies in Loops
  - Data Movement

The compiler needs your higher level perspective (in the form of directive hints) to get correct results and reasonable performance.



#### **Data Dependencies**

Most directive based parallelization consists of splitting up big do/for loops into independent chunks that the many processors can work on simultaneously.

Take, for example, a simple for loop like this:

for(index=0; index<1000000; index++)
 Array[index] = 4 \* Array[index];</pre>

When run on 1000 processors, it will execute something like this...



## No Data Dependency





#### Data Dependency

But what if the loops are not entirely independent?

Take, for example, a similar loop like this:

#### 

This is perfectly valid serial code.



#### Data Dependency

Now Processor 1, in trying to calculate its first iteration...

needs the result of Processor O's last iteration. If we want the correct ("same as serial") result, we need to wait until processor O finishes. Likewise for processors 2, 3, ...



#### **Data Dependencies**

That is a data dependency. If the compiler even <u>suspects</u> that there is a data dependency, it will, for the sake of correctness, refuse to parallelize that loop with kernels.

11, Loop carried dependence of 'Array' prevents parallelization Loop carried backward dependence of 'Array' prevents vectorization

As large, complex loops are quite common in HPC, especially around the most important parts of your code, the compiler will often balk most when you most need a kernel to be generated. What can you do?



## **Data Dependencies**

Rearrange your code to make it more obvious to the compiler that there is not really a data dependency.

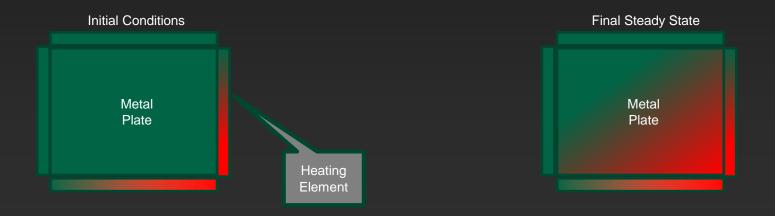
Eliminate a real dependency by changing your code.

- There is a common bag of tricks developed for this as this issue goes back 40 years in HPC. Many are quite trivial to apply.
- The compilers have gradually been learning these themselves.
- Override the compiler's judgment (independent clause) at the risk of invalid results. Misuse of restrict has similar consequences.



#### **Our Foundation Exercise: Laplace Solver**

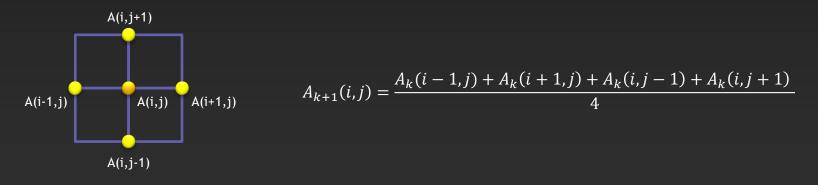
- I've been using this for MPI, OpenMP and now OpenACC. It is a great simulation problem, not rigged for OpenACC.
- In this most basic form, it solves the Laplace equation:  $abla^2 f(x,y) = oldsymbol{0}$
- The Laplace Equation applies to many physical problems, including:
  - Electrostatics
  - Fluid Flow
  - Temperature
- For temperature, it is the Steady State Heat Equation:





#### **Exercise Foundation: Jacobi Iteration**

- The Laplace equation on a grid states that each grid point is the average of it's neighbors.
- We can iteratively converge to that state by repeatedly computing new values at each point from the average of neighboring points.
- We just keep doing this until the difference from one pass to the next is small enough for us to tolerate.





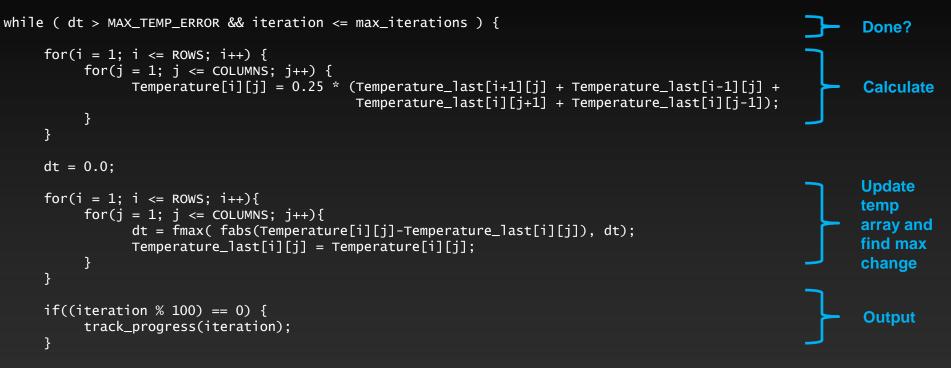
## Serial Code Implementation







# Serial C Code (kernel)



iteration++;



## Serial C Code Subroutines

#### void initialize(){

```
int i,j;
for(i = 0; i <= ROWS+1; i++){</pre>
    for (j = 0; j \le COLUMNS+1; j++)
        Temperature_last[i][j] = 0.0;
}
// these boundary conditions never change throughout run
// set left side to 0 and right to a linear increase
for(i = 0; i <= ROWS+1; i++) {</pre>
    Temperature_last[i][0] = 0.0;
    Temperature_last[i][COLUMNS+1] = (100.0/ROWS)*i;
}
// set top to 0 and bottom to linear increase
for(j = 0; j <= COLUMNS+1; j++) {</pre>
    Temperature_last[0][j] = 0.0;
    Temperature_last[ROWS+1][j] = (100.0/COLUMNS)*j;
```

BCs could run from 0 to ROWS+1 or from 1 to ROWS. We chose the former.

```
void track_progress(int iteration) {
    int i;
    printf("-- Iteration: %d --\n", iteration);
    for(i = ROWS-5; i <= ROWS; i++) {
        printf("[%d,%d]: %5.2f ", i, i,Temperature[i][i]);
    }
    printf("\n");
}</pre>
```



#### Whole C Code

#include <stdlib.h>
#include <stdio.h>
#include <math.h>
#include <sys/time.h>

// size of plate #define COLUMNS 1000 #define ROWS 1000

// largest permitted change in temp (This value takes about 3400 steps)
#define MAX\_TEMP\_ERROR 0.01

double Temperature[ROWS+2][COLUMNS+2]; // temperature grid double Temperature\_last[ROWS+2][COLUMNS+2]; // temperature grid from last iteration

// helper routines
void initialize();
void track\_progress(int iter);

int main(int argc, char \*argv[]) {

int i, j; // grid indexes
int max\_iterations; // number of iterations
int iteration=1; // current iteration
double dt=100; tr\_time, stop\_time, elapsed\_time; // timers

printf("Maximum iterations [100-4000]?\n"); scanf("%d", &max\_iterations);

gettimeofday(&start\_time,NULL); // Unix timer

initialize();

// initialize Temp\_last including boundary conditions

```
// do until error is minimal or until max steps
while ( dt > MAX_TEMP_ERROR && iteration <= max_iterations ) {</pre>
```

```
// main calculation: average my four neighbors
for(i = 1; i <= ROWS; i++) {
   for(j = 1; j <= COLUNNS; j++) {
      Temperature[i][j] = 0.25 * (Temperature_last[i+1][j] + Temperature_last[i][j+1] + Temperature_last[i][j-1]);
      }
}</pre>
```

```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
for(i = 1; i <= ROWS; i++){
  for(j = 1; j <= COLUMNS; j++){
    dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
    Temperature_last[i][j] = Temperature[i][j];
  }
}
// periodically print test values
```

```
// periodically print test values
if((iteration % 100) == 0) {
    track_progress(iteration);
}
```

iteration++;

gettimeofday(&stop\_time,NULL); timersub(&stop\_time, &start\_time, &elapsed\_time); // Unix time subtract routine

printf("\nMax error at iteration %d was %f\n", iteration-1, dt); printf("Total time was %f seconds.\n", elapsed\_time.tv\_sec+elapsed\_time.tv\_usec/1000000.0);

```
// initialize plate and boundary conditions
// Temp_last is used to to start first iteration
void initialize(){
```

```
for(i = 0; i <= ROWS+1; i++){
    for(j = 0; j <= COLUMNS+1; j++){
        Temperature_last[i][j] = 0.0;
    }
}</pre>
```

// these boundary conditions never change throughout run

```
// set left side to 0 and right to a linear increase
for(i = 0; i <= ROWS+1; i++) {
    Temperature_last[i][0] = 0.0;
    Temperature_last[i][COLUMNS+1] = (100.0/ROWS)*i;
}
```

```
// set top to 0 and bottom to linear increase
for(j = 0; j <= COLUMNS+1; j++) {
    Temperature_last[0][j] = 0.0;
    Temperature_last[ROWS+1][j] = (100.0/COLUMNS)*j;
```

```
ر
```

3

int i.j:

// print diagonal in bottom right corner where most action is void track\_progress(int iteration)  $\{$ 

int i;

```
printf("------ Iteration number: %d ------\n", iteration);
for(i = ROWS-5; i <= ROWS; i++) {
    printf("[%d,%d]: %5.2f ", i, i, Temperature[i][i]);
}
printf("\n");
```



# Serial Fortran Code (kernel)

```
do while ( dt > max_temp_error .and. iteration <= max_iterations)
                                                                                                      Done?
  do j=1,columns
    do i=1.rows
        temperature(i,j)=0.25*(temperature_last(i+1,j)+temperature_last(i-1,j)+ &
                                                                                                      Calculate
                               temperature_last(i,j+1)+temperature_last(i,j-1) )
    enddo
  enddo
 dt=0.0
                                                                                                      Update
  do j=1,columns
    do i=1, rows
                                                                                                      temp
        dt = max( abs(temperature(i,j) - temperature_last(i,j)), dt )
                                                                                                      array and
        temperature_last(i,j) = temperature(i,j)
                                                                                                      find max
    enddo
                                                                                                      change
  enddo
  if( mod(iteration, 100).eq.0 ) then
                                                                                                      Output
    call track_progress(temperature, iteration)
  endif
  iteration = iteration+1
```

#### enddo



#### Serial Fortran Code Subroutines

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,j

double precision, dimension(0:rows+1,0:columns+1) :: temperature\_last

 $temperature_last = 0.0$ 

!these boundary conditions never change throughout run

!set left side to 0 and right to linear increase do i=0,rows+1 temperature\_last(i,0) = 0.0 temperature\_last(i,columns+1) = (100.0/rows) \* i enddo

```
!set top to 0 and bottom to linear increase
do j=0,columns+1
    temperature_last(0,j) = 0.0
    temperature_last(rows+1,j) = ((100.0)/columns) * j
enddo
```

end subroutine initialize

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,iteration

double precision, dimension(0:rows+1,0:columns+1) :: temperature



#### Whole Fortran Code

#### program serial implicit none

!Size of plate integer, parameter :: columns=1000 integer, parameter :: rows=1000 double precision, parameter :: max\_temp\_error=0.01

integer double precision real :: i, j, max\_iterations, iteration=1
:: dt=100.0
:: start\_time, stop\_time

double precision, dimension(0:rows+1,0:columns+1) :: temperature, temperature\_last

print\*, 'Maximum iterations [100-4000]?"
read\*, max\_iterations

call initialize(temperature\_last)

!do until error is minimal or until maximum steps
do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

#### do j=1,columns

do i=1,rows
 temperature(i,j)=0.25\*(temperature\_last(i+1,j)+temperature\_last(i-1,j)+ &
 temperature\_last(i,j+1)+temperature\_last(i,j-1) )

```
enddo
enddo
```

dt=0.0

!copy grid to old grid for next iteration and find max change do j=1,columns do i=1,rows dt = max( abs(temperature(i,j) - temperature\_last(i,j)), dt ) temperature\_last(i,j) = temperature(i,j) enddo enddo

```
!periodically print test values
if( mod(iteration,100).eq.0 ) then
    call track_progress(temperature, iteration)
endif
```

#### iteration = iteration+1

#### enddo

call cpu\_time(stop\_time)

print\*, 'Max error at iteration ', iteration-1, ' was ',dt
print\*, 'Total time was ',stop\_time-start\_time, ' seconds.'

end program serial

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,j

double precision, dimension(0:rows+1,0:columns+1) :: temperature\_last

 $temperature_last = 0.0$ 

!these boundary conditions never change throughout run

!set left side to 0 and right to linear increase do i=0,rows+1 temperature\_last(i,0) = 0.0 temperature\_last(i,columns+1) = (100.0/rows) \* i enddo

!set top to 0 and bottom to linear increase do j=0,columns+1 temperature\_last(0,j) = 0.0 temperature\_last(rows+1,j) = ((100.0)/columns) \* j enddo

end subroutine initialize

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,iteration

double precision, dimension(0:rows+1,0:columns+1) :: temperature



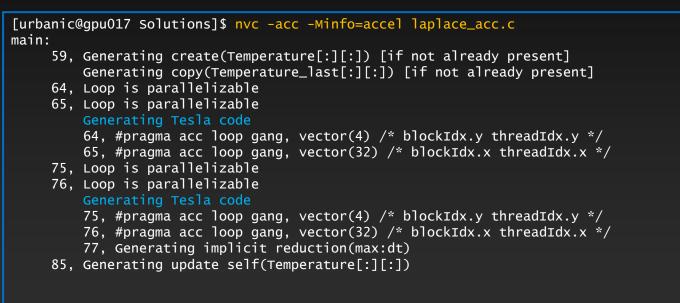
### **Exercises: General Instructions for Compiling**

- Exercises are in the "Exercises/OpenACC" directory in your home directory
- Solutions are in the "Solutions" subdirectory
- To compile nvc -acc laplace.c nvfortran -acc laplace.f90
- This will generate the executable a.out



# Exercises: Very useful compiler option

Adding -Minfo=accel to your compile command will give you some very useful information about how well the compiler was able to honor your OpenACC directives.





# Special Instructions for Running on the GPUs (during this workshop)

As mentioned, on Bridges2 you generally only have to use the queueing system when you want to. However, as we have hundreds of you wanting quick turnaround, we will have to use it today.

Once you have an a.out that you want to run, you should use the simple job that we have already created (in Exercises/OpenACC) for you to run:

fred@bridges2-login011\$ sbatch gpu.job



#### **Output From Your Batch Job**

The machine will tell you it submitted a batch job, and you can await your output, while will come back in a file with the corresponding number as a name:

slurm-138555.out

As everything we are doing this afternoon only requires a few minutes at most (and usually just seconds), you could just sit there and wait for the file to magically appear. At which point you can "more" it or review it with your editor.



# **Changing Things Up**

If you get impatient, or want to see what the machine us up to, you can look at the situation with squeue.

You might wonder what happened to the interaction count that the user is prompted for. I stuck a reasonable default (4000 iterations) into the job file. You can edit it if you want to. The whole job file is just a few lines.

Congratulations, you are now a Batch System veteran. Welcome to supercomputing.



# Exercise 1: Using kernels to parallelize the main loops (About 20 minutes)

Q: Can you get a speedup with just the kernels directives?

#### 1. Edit *laplace\_serial.c/f*90

- 1. Maybe copy your intended OpenACC version to *laplace\_acc.c* to start
- 2. Add directives where it helps
- 2. Compile with OpenACC parallelization
  - 1. nvc -acc -Minfo=accel laplace\_acc.c or nvfortran -acc -Minfo=accel laplace\_acc.f90
  - 2. Look at your compiler output to make sure you are having an effect
- 3. Run
  - 1. sbatch gpu.job (Leave it at 4000 iterations if you want a solution that converges to current tolerance)
  - 2. Look at output in file that returns (something like slurm-138555.out)
  - 3. Compare the serial and your OpenACC version for performance difference



### **Exercise 1 C Solution**

while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>

}

```
Generate a GPU kernel
#pragma acc kernels
for(i = 1; i <= ROWS; i++) {</pre>
    for(j = 1; j <= COLUMNS; j++) {</pre>
        Temperature[i][j] = 0.25 * (Temperature_last[i+1][j] + Temperature_last[i-1][j] +
                                      Temperature_last[i][j+1] + Temperature_last[i][j-1]);
    }
}
dt = 0.0; // reset largest temperature change
                                                                                      Generate a GPU kernel
#pragma acc kernels
for(i = 1; i <= ROWS; i++){</pre>
    for(j = 1; j <= COLUMNS; j++){</pre>
        dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
        Temperature_last[i][i] = Temperature[i][i];
    }
}
if((iteration % 100) == 0) {
    track_progress(iteration);
}
iteration++;
```

### **Exercise 1 Fortran Solution**

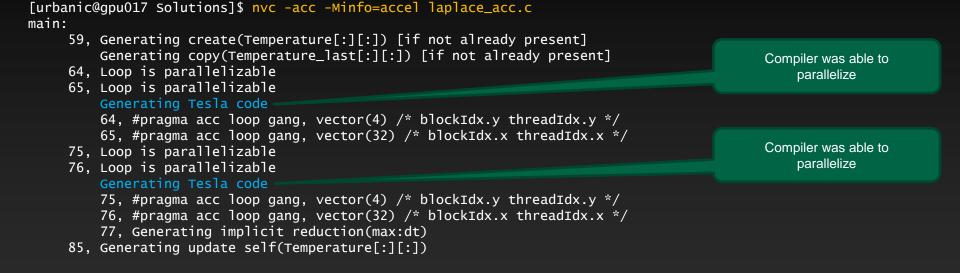
do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

enddo

```
Generate a GPU kernel
!$acc kernels
do j=1,columns
   do i=1, rows
      temperature(i,j)=0.25*(temperature_last(i+1,j)+temperature_last(i-1,j)+ &
                             temperature_last(i,j+1)+temperature_last(i,j-1) )
   enddo
enddo
!$acc end kernels
dt=0.0
                                                                                    Generate a GPU kernel
!$acc kernels
do j=1,columns
   do i=1,rows
      dt = max( abs(temperature(i,j) - temperature_last(i,j)), dt )
      temperature_last(i,j) = temperature(i,j)
   enddo
enddo
!$acc end kernels
if( mod(iteration, 100).eq.0 ) then
   call track_progress(temperature, iteration)
endif
iteration = iteration+1
```

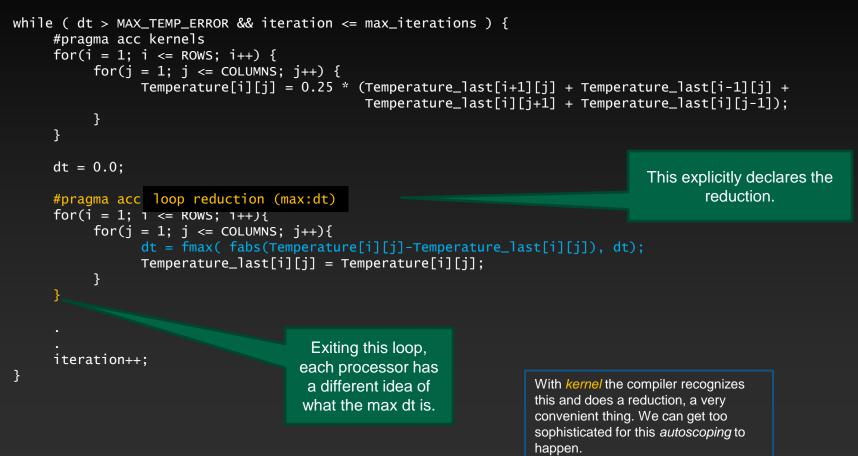


# Exercise 1: Compiler output (C)





#### First, about that "reduction"





# Exercise 1: Performance

#### 3372 steps to convergence

Execution	Time (s)	Speedup
CPU Serial	20.6	
CPU 2 OpenMP threads	10.3	2.0
CPU 4 OpenMP threads	5.2	4.0
CPU 8 OpenMP threads	2.6	7.9
CPU 16 OpenMP threads	1.4	14.7
CPU 32 OpenMP threads	0.80	25.7
CPU 64 OpenMP threads	0.72	28.6
CPU 128 OpenMP threads	1.4	14.7
OpenACC GPU	32.4	0.6x



Using NVHPC 21.2 on a V100

# What's with the OpenMP?

We can compare our GPU results to the <u>best</u> the multi-core CPUs can do.

If you are familiar with OpenMP, or even if you are not, you can compile and run the OpenMP enabled versions in your OpenMP directory as:

nvc -mp laplace\_omp.c or nvfortran -mp laplace\_omp.f90

then to run on 8 threads do:

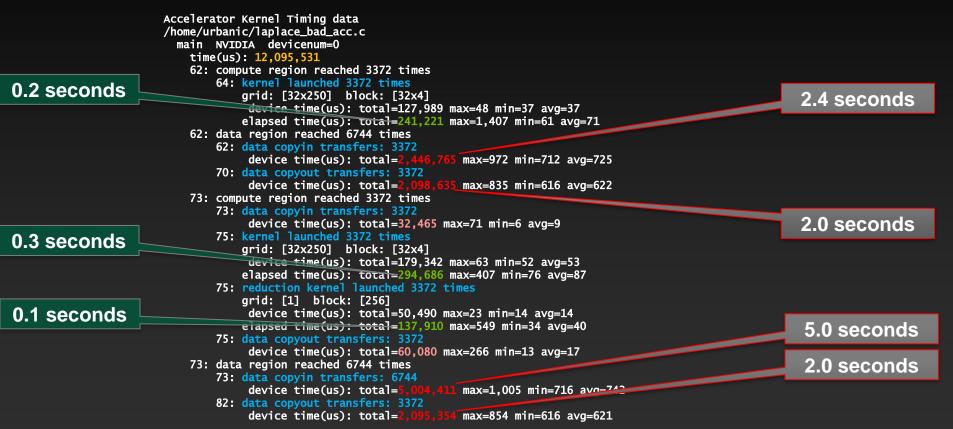
```
export OMP_NUM_THREADS=8
a.out
```

Note that you probably only have 8 real cores if you are still on a GPU node. Do something like "interact -n28" if you want a full node of cores.



#### What went wrong?

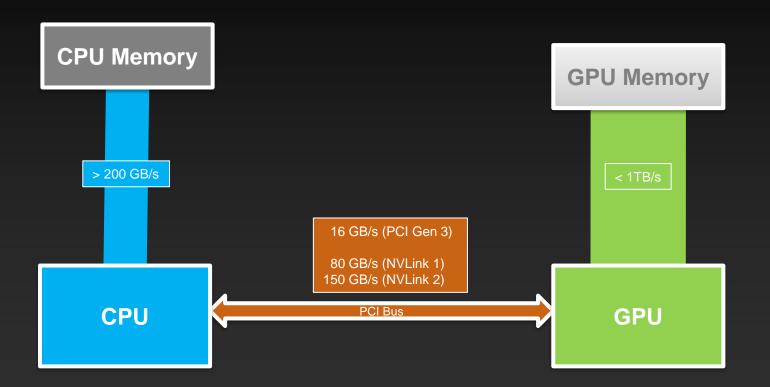
#### export PGI\_ACC\_TIME=1 to activate profiling and run again:





### **Basic Concept**

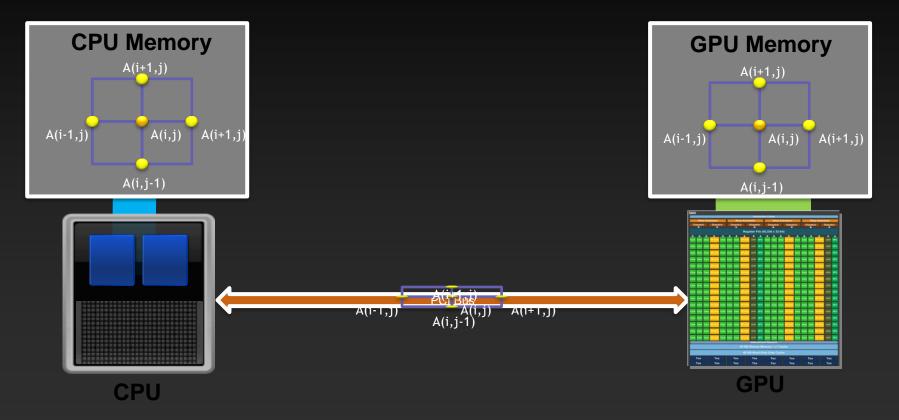
Simplified, but sadly true





All bandwidths one-direction.

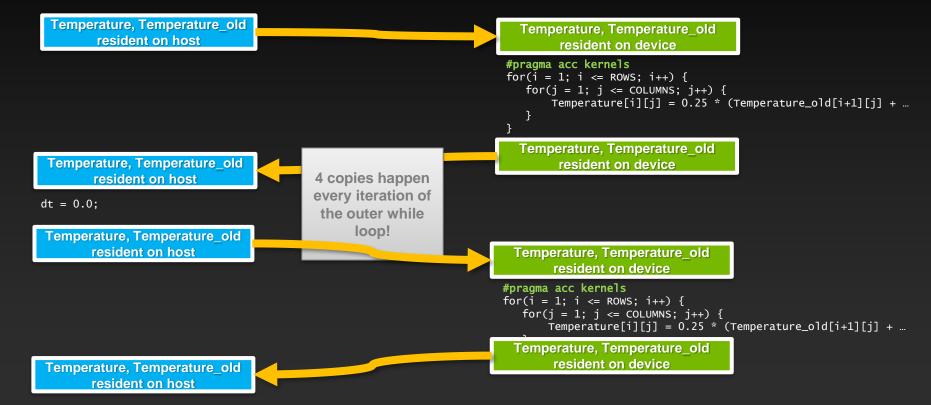
#### **Multiple Times Each Iteration**





### **Excessive Data Transfers**

#### while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>





#### Data Management

The First, Most Important, and Possibly Only OpenACC Optimization



### Scoped Data Construct Syntax

#### Fortran

С

!\$acc data [clause ...]
 structured block
!\$acc end data

#pragma acc data [clause ...]
{
 structured block
}



#### **Data Clauses**

copy( list )

Allocates memory on GPU and copies data from host to GPU when entering region and copies data to the host when exiting region. Principal use: For many important data structures in your code, this is a logical default to input, modify and return the data.

copyin( list ) Allocates memory on GPU and copies data from host to GPU when
entering region.
Principal use: Think of this like an array that you would use as just
an input to a subroutine.

copyout( list ) Allocates memory on GPU and copies data to the host when exiting
region.
Principal use: A result that isn't overwriting the input data structure.

create( list ) Allocates memory on GPU but does not copy.
Principal use: Temporary arrays.



# Array Shaping

Compilers sometimes cannot determine the size of arrays, so we must specify explicitly using data clauses with an array "shape". The compiler will let you know if you need to do this. Sometimes, you will want to for your own efficiency reasons.

#pragma acc data copyin(a[0:size]), copyout(b[s/4:3\*s/4])

Fortran

C

!\$acc data copyin(a(1:size)), copyout(b(s/4:3\*s/4))

- Fortran uses start:end and C uses start:length
- Data clauses can be used on data, kernels or parallel



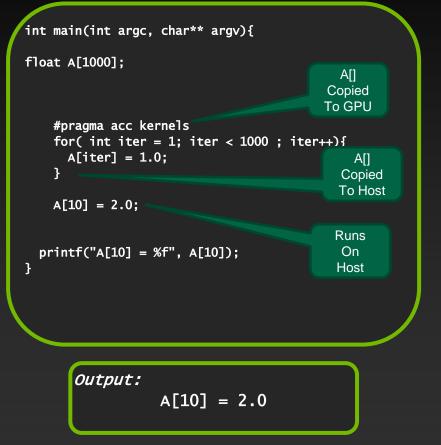
#### Compiler will (increasingly) often make a good guess...

```
int main(int argc, char *argv[]) {
                                                        nvc -acc -Minfo=accel loops.c
                                                        main:
 int i;
                                                            6, Generating present_or_copyout(C[:])
 double A[2000], B[1000], C[1000];
                                                               Generating present_or_copy(B[:])
                                                               Generating present_or_copyout(A[:1000])
                                          Smarter
                                                               Generating NVIDIA code
                                                            7, Loop is parallelizable
 #pragma acc kernels
                                         Smartest
                                                               Accelerator kernel generated
 for (i=0; i<1000; i++){
   A[i] = 4 * i;
   B[i] = B[i] + 2;
   C[i] = A[i] + 2 * B[i];
```

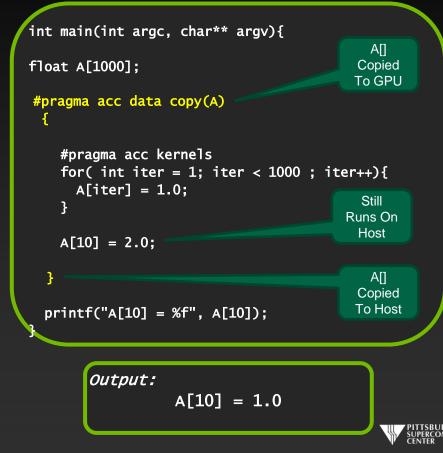


#### **Data Regions Have Real Consequences**

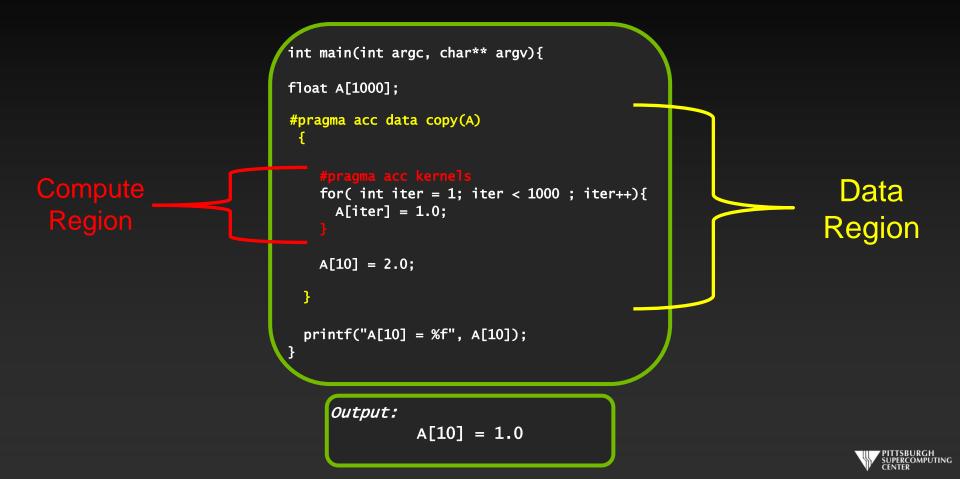
#### Simplest Kernel



#### With Global Data Region



#### Data Regions Are Different Than Compute Regions



#### **Data Movement Decisions**

- Much like loop data dependencies, sometime the compiler needs your human intelligence to make high-level decisions about data movement. Otherwise, it must remain conservative - sometimes at great cost.
- You must think about when data truly needs to migrate, and see if that is better than the default.
- Besides the scope-based data clauses, there are OpenACC options to let us manage data movement more intensely or asynchronously. We could manage the above behavior with the update construct:

Fortran :
!\$acc update [host(), device(), ...]

C: #pragma acc update [host(), device(), ...]

Ex: #pragma acc update host(Temp\_array) //Get host a copy from device



# Exercise 2: Use acc data to minimize transfers

(about 40 minutes)

Q: What speedup can you get with data + kernels directives?

- Start with your Exercise 1 solution or grab laplace\_bad\_acc.c/f90 from the Solutions subdirectory. This is just the solution of the last exercise.
- Add data directives where it helps.
  - Think: when *should* I move data between host and GPU? Think how you would do it by hand, then determine which data clauses will implement that plan.
  - Hint: you may find it helpful to ignore the output at first and just concentrate on getting the solution to converge quickly (at 3372 steps). Then worry about *updating* the printout.



#### Exercise 2 C Solution

```
#pragma acc data copy(Temperature_last, Temperature)
while ( dt > MAX_TEMP_ERROR && iteration <= max_iterations ) {</pre>
```

No data movement in this block.

```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
#pragma acc kernels
for(i = 1; i <= ROWS; i++){
    for(j = 1; j <= COLUMNS; j++){
        dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
        Temperature_last[i][j] = Temperature[i][j];
    }
}</pre>
```

```
// periodically print test values
if((iteration % 100) == 0) {
    #pragma acc update host(Temperature)
    track_progress(iteration);
}
```

Except once in a while here.



#### Exercise 2, Slightly better solution

#pragma acc data copy(Temperature\_last), create(Temperature)
while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>



```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
#pragma acc kernels
for(i = 1; i <= ROWS; i++){
   for(j = 1; j <= COLUMNS; j++){
     dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
     Temperature_last[i][j] = Temperature[i][j];
   }
}</pre>
```

```
// periodically print test values
if((iteration % 100) == 0) {
    #pragma acc update host(Temperature)
    track_progress(iteration);
}
```

iteration++;



#### Slightly better still solution

```
#pragma acc data copy(Temperature_last), create(Temperature)
while ( dt > MAX_TEMP_ERROR && iteration <= max_iterations ) {</pre>
```

```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
#pragma acc kernels
for(i = 1; i <= ROWS; i++){
    for(j = 1; j <= COLUMNS; j++){
        dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
        Temperature_last[i][j] = Temperature[i][j];
    }
}</pre>
```

```
// periodically print test values
if((iteration % 100) == 0) {
    #pragma acc update host(Temperature[ROWS-4:5][COLUMNS-4:5])
    track_progress(iteration);
}
```



Only need corner elements.



#### **Exercise 2 Fortran Solution**

!\$acc data copy(temperature\_last), create(temperature)
do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

!\$acc end data

```
!$acc kernels
  do i=1.columns
     do i=1, rows
        temperature(i,j)=0.25*(temperature_last(i+1,j)+temperature_last(i-1,j)+ &
                                temperature_last(i,j+1)+temperature_last(i,j-1) )
     enddo
  enddo
   !$acc end kernels
  dt=0.0
   !copy grid to old grid for next iteration and find max change
   !$acc kernels
  do j=1.columns
     do i=1.rows
        dt = max( abs(temperature(i,j) - temperature_last(i,j)), dt )
        temperature_last(i,j) = temperature(i,j)
     enddo
  enddo
                                                           !$acc update host(temperature(columns-5:columns,rows-5:rows))
   !$acc end kernels
   !periodically print test values
                                                                                                        Except bring back a copy
  if( mod(iteration, 100).eq.0 ) then
                                                                                                                     here
      !$acc update host(temperature)
     call track_progress(temperature, iteration)
  endif
  iteration = iteration+1
enddo
```

Keep these on GPU

# Exercise 2: Performance

#### 3372 steps to convergence

Execution	Time (s)	Speedup
CPU Serial	20.6	
CPU 2 OpenMP threads	10.3	2.0
CPU 4 OpenMP threads	5.2	4.0
CPU 8 OpenMP threads	2.6	7.9
CPU 16 OpenMP threads	1.4	14.7
CPU 32 OpenMP threads	0.80	25.7
CPU 64 OpenMP threads	0.72	28.6
CPU 128 OpenMP threads	1.4	14.7
OpenACC GPU	0.62	33



# **OpenACC or OpenMP?**

Don't draw any grand conclusions yet. We have gotten impressive speedups from both approaches. But our problem size is pretty small. Our main data structure is:

1000 x 1000 = 1M elements = 8MB of memory

We have 2 of these (temperature and temperature\_last) so we are using roughly 16 MB of memory. Not very large. When divided over cores it gets even smaller and can easily fit into cache.

The algorithm is realistic, but the problem size is tiny and hence the memory bandwidth stress is very low.



# **OpenACC or OpenMP on Larger Data?**

We can easily scale this problem up, so why don't I? Because it is nice to have exercises that finish in a few minutes or less.

We scale this up to  $10K \times 10K$  (1.6 GB problem size) for the hybrid challenge. These numbers start to look a little more realistic. But the serial code takes over 30 minutes to finish. That would have gotten us off to a slow start!

Execution	Time (s)	Speedup
CPU Serial	2187	
CPU 16 OpenMP threads	183	12
CPU 28 OpenMP threads	162	13.5
OpenACC	103	21

Obvious cusp for core scaling appears

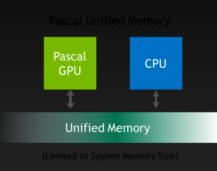
#### 10K x 10K Problem Size



#### Latest Happenings In Data Management

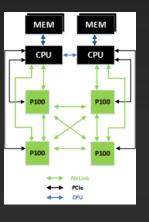
#### Unified Memory

- Unified address space allows us to pretend we have shared memory
- Skip data management, hope it works, and then optimize if necessary
- For dynamically allocated memory can eliminate need for pointer clauses



#### NVLink

One route around PCI bus (with multiple GPUs)





### **Further speedups**

OpenACC gives us even more detailed control over parallelization • Via gang, worker, and vector clauses

By understanding more about OpenACC execution model and GPU hardware organization, we can get higher speedups on this code

By understanding bottlenecks in the code via profiling, we can reorganize the code for higher performance

 But you have already gained most of any potential speedup, and you did it with a few lines of directives!



# Is OpenACC Living Up To My Claims?

- High-level. No involvement of OpenCL, CUDA, etc.
- Single source. No forking off a separate GPU code. Compile the same program for accelerators or serial; non-GPU programmers can play along.
- Efficient. Experience shows very favorable comparison to low-level implementations of same algorithms. kernels is magical!
- Performance portable. Supports GPU accelerators and co-processors from multiple vendors, current and future versions.
- Incremental. Developers can port and tune parts of their application as resources and profiling dictates. No wholesale rewrite required. Which can be <u>quick.</u>



#### In Conclusion...

